

# Algorithm verification: Sample problems

The problem below, with solution, is a model to help solving problems in assignment 2.

## 1. Sum problem

Find the sum of the elements on an array of numbers.

### Iterative solution with assertions:

#### Sum(A)

> Precondition: A is an array of numbers

$i \leftarrow 1$

$y \leftarrow 0$

while  $i \leq |A|$

    > Loop invariant:  $y = A[1] + A[2] + \dots + A[i - 1]$

$y \leftarrow y + A[i]$

$i \leftarrow i + 1$

return y

> Postcondition:  $y = A[1] + A[2] + \dots + A[|A|]$

### Correctness proof:

We proceed by induction.

1. (Base) The loop invariant holds on the first iteration, because y is 0 and  $A[i - 1]$  refers to no element of A.
2. (Induction) If the loop invariant holds on the  $n$ th iteration, then it will hold on the  $(n + 1)$ st, because on each iteration  $A[i]$  is added to y, and the element  $A[i]$  takes the name  $A[i - 1]$  after the incrementing of  $i$ .
3. The precondition follows from the loop invariant, because at the beginning of the last iteration,  $i = |A| - 1$ , and that iteration adds 1 to  $i$  and adds  $A[|A|]$  to y.
4. Hence at the end of the last iteration,  $y = A[1] + A[2] + \dots + A[i]$ , which is the same assertion as the postcondition, since  $i = |A|$ .

## 2. Last-zero problem

*Problem:* Find the location of the rightmost zero in a bit array.

*Solution:*

#### Last0(A)

// Pre:  $|A| > 0$  and A contains a 0

$i \leftarrow 1$

while  $i \leq |A|$

> Loop invariant: y stores location of

> rightmost 0 in  $A[1 .. i - 1]$

    If  $A[i] = 0$

$y \leftarrow i$

$i \leftarrow i + 1$

return y

// Post: y stores location of rightmost 0 in A

### Correctness proof:

We proceed by induction. [To be completed.]