# Godel's Theorem: Limits of logic and computation

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# **Overview**

- Kurt Gödel, 1931, at age 25, in Vienna, shook up the foundations of mathematical thought as developed in Russell and Whitehead's *Principia Mathematica*
- He showed that no formal axiomatic system of logic could be both *consistent* and *complete*
- Since predicate calculus was proven consistent, Gödel's theorem showed that some truths are unprovable in the system

## **General-purpose definitions**

• *Consistency:* Attribute of a system in which no false assertion can be proven

 $\forall p(|-p \Rightarrow p)$ 

• *Completeness:* Attribute of a system in which every true assertion can be proven

 $\forall p(p \Rightarrow |-p)$ 

# Intuition of Gödel's counter-example

• Gödel showed that *something like* the following is true but unprovable:

"This assertion has no proof"

• But the assertion actually used has a precise meaning

#### **Definitions related to theorem**

- *Proof(x, y, z):* The sequence of symbols with Gödel number *x* is a valid proof that  $S_y(z)$  is true, where  $S_y$  is the formula with Gödel number *y*, and *z* is a natural number.
- Let G(y) be the assertion that  $\neg \exists x \operatorname{Proof}(x, y, y)$ i.e., G(y) means that what the formula with Gödel number y asserts about y is unprovable
- *Example*: Where y is the predicate *even*, G(y) is the (false) assertion that one cannot prove that the Gödel number of the predicate *even* is an even number

# Gödel's theorem

 $\neg \exists x \operatorname{Proof}(x, \#(G), \#(G)) \text{ is true in}$ predicate calculus, but cannot be proven in predicate calculus

- *Or:* What *G* asserts about predicate *G* is true but unprovable
- *Or:* There is no proof that *G*(#(*G*)) is true, but *G*(#(*G*)) is nevertheless true

#### **Proof of Gödel's theorem**

- 1. Suppose the theorem's claim is false and there exists a proof in predicate calculus that  $\neg \exists x \operatorname{Proof}(x, \#(G), \#(G))$ .
- 2. Then the claim is true, by consistency of predicate calculus.
- 3. Therefore by its own assertion it is unprovable in predicate calculus.
- 4. This contradicts step 1, hence claim holds.

*Note:* The above proof goes outside the formal system of predicate calculus.

# What this has to do with computer science

- Computations and sets of computations are proofs about programs
- Gödel's theorem was precursor to Turing's result on the Halting Problem

## Larger significance

- Gödel's theorem was part of a series of results establishing limits of knowledge that also included
  - Einstein's relativity results
  - Heisenberg's Uncertainty Principle
- It marked the failure of Russell's, Whitehead's and Hilbert's effort to reduce mathematics to first-order logic and set theory

#### Source

• Flake, *The computational beauty of nature*, MIT, 2001.

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