# Arithmetical Structures on Graphs 

## Darren Glass

## Gettysburg College

(Joint work with various folks)

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$$
M=\left[\begin{array}{cccc}
d_{1} & -1 & -1 & -1 \\
-1 & d_{2} & -1 & 0 \\
-1 & -1 & d_{3} & -1 \\
-1 & 0 & -1 & d_{4}
\end{array}\right]
$$

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\end{array}\right]
$$

If there exists a vector $\mathbf{r}=\left(r_{1}, \ldots, r_{n}\right)$ so that:

- All $r_{i}$ are positive integers
- $\operatorname{gcd}\left(r_{1}, \ldots, r_{n}\right)=1$
- $\mathbf{r} \cdot M=\overrightarrow{0}$
then we say that $(\mathbf{r}, \mathbf{d})$ is an arithmetical structure on $G$.

Note that for any d, $M$ will have rank $n-1$ so $\mathbf{r}$ is unique if it exists.

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then we say that $(\mathbf{r}, \mathbf{d})$ is an arithmetical structure on $G$.

Note that for any $\mathbf{d}, M$ will have rank $n-1$ so $\mathbf{r}$ is unique if it exists. (and vice versa)

## Example



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$$
M=\left[\begin{array}{cccc}
3 & -1 & -1 & -1 \\
-1 & 2 & -1 & 0 \\
-1 & -1 & 3 & -1 \\
-1 & 0 & -1 & 2
\end{array}\right] \text { is the Laplacian of the graph. }
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$M=\left[\begin{array}{cccc}3 & -1 & -1 & -1 \\ -1 & 2 & -1 & 0 \\ -1 & -1 & 3 & -1 \\ -1 & 0 & -1 & 2\end{array}\right]$ is the Laplacian of the graph.
All rows sum to zero $\Rightarrow(1,1, \ldots, 1) \cdot M=\overrightarrow{0}$

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All rows sum to zero $\Rightarrow(1,1, \ldots, 1) \cdot M=\overrightarrow{0}$
$((1,1, \ldots, 1), \mathbf{d})$ is an arithmetical structure on $G$.

## Parenthetical Aside

For all graphs $G$, we have that $(L,(1, \ldots, 1))$ is an arithmetical structure on $G$.

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## Definition

Let $\mathcal{D}=\left\{\left(d_{1}, \ldots, d_{n}\right) \mid \sum d_{i}=0\right\}$. The Jacobian of $G$ is defined to be $\mathcal{D} / \operatorname{Im}(L)$.

Analogous to the Jacobian of a curve in algebraic geometry, and studied by many people (Bak, Baker-Norine, Propp, ...) under many different names (Critical Groups, Sandpile Groups, ...).

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For any arithmetical structure ( $\mathbf{r}, \mathbf{d}$ ) on $G$ we can define an analogue to the Jacobian by considering $\operatorname{Ker}(\mathbf{r}) / \operatorname{Im}(M)$. Can think of this as extending the analogy to curves with components of higher multiplicity.

## Another Example



Let $M=\left[\begin{array}{cccc}3 & -1 & -1 & -1 \\ -1 & 1 & -1 & 0 \\ -1 & -1 & 11 & -1 \\ -1 & 0 & -1 & 1\end{array}\right]$ and $\mathbf{r}=(3,4,1,4)$.
Then $\mathbf{r} \cdot M=\overrightarrow{0}$
$\Rightarrow(\mathbf{r},(3,1,11,1))$ is an arithmetical structure on $G$.

## What Does This Have To Do With Number Theory?

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## Definition

An arithmetical structure on $G$ is a labeling of the vertices of $G$ with nonnegative integers $r_{i}$ so that $\operatorname{gcd}\left(r_{1}, \ldots, r_{n}\right)=1$ and for each vertex $i$ we have that $r_{i} \mid \sum_{(i, j) \in E(G)} r_{j}$


An arithmetical structure is a vector $\left(r_{1}, r_{2}, r_{3}, r_{4}\right)$ so that:

- $r_{1} \mid r_{2}+r_{3}+r_{4}$
- $r_{2} \mid r_{1}+r_{3}$
- $r_{3} \mid r_{1}+r_{2}+r_{4}$
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## Question

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For a fixed graph, there are finitely many arithmetical structures.

## For small graphs, use Mathematica!

```
In[88]= Do[
        Do[Do[Do[If[GCD[a,b, c, d] == 1, If [Divisible[b+c+d,a], If[Divisible[a + c, b], If[
            Divisible[a+c, d], If[Divisible[a+b+d, c], AppendTo[S,{a,b, c, d}]]]]]],
                {a,1,50}], {b, 1, 50}], {c, 1, 50}], {d, 1, 50}]
In[89]= S
Ou{[89]= {{1, 1, 1, 1}, {3,1, 1, 1}, {1, 2, 1, 1}, {5, 3, 1, 1}, {3, 4, 1, 1}, {2, 1, 2, 1},
    {4, 1, 2, 1}, {8, 5, 2, 1}, {1, 1, 3, 1}, {3, 2, 3, 1}, {1, 4, 3, 1}, {7, 10, 3, 1},
    {2, 1, 4, 1}, {6, 1, 4, 1}, {8, 3, 4, 1}, {14, 9, 4, 1}, {1, 3, 5, 1}, {4, 1, 6, 1},
    {3,10, 7, 1}, {4, 3, 8, 1}, {2, 5, 8, 1}, {18,5, 12, 1}, {4, 9, 14, 1},
    {12, 5, 18, 1}, {1, 1, 1, 2}, {1, 2, 1, 2}, {5, 2, 1, 2}, {7, 4, 1, 2}, {3, 1, 3, 2},
    {9,4,3,2}, {1, 2, 5, 2}, {1, 4, 7, 2}, {3,4, 9, 2}, {5, 1, 1, 3}, {5,6,1,3},
    {4,3,2,3}, {8, 1, 4, 3}, {2,3,4,3}, {1, 1, 5, 3}, {1, 6, 5, 3}, {4, 1, 8, 3},
    {3,1,1,4}, {7, 2, 1, 4}, {3,4,1,4}, {11,6,1, 4}, {1, 1, 3, 4}, {9, 2, 3, 4},
    {1,4,3,4}, {1, 2, 7, 4}, {3, 2, 9, 4}, {1,6,11, 4}, {8, 1, 2, 5}, {2, 1, 8, 5},
    {18, 1, 12, 5}, {12, 1, 18, 5}, {5, 3, 1, 6}, {11, 4, 1, 6}, {1, 3, 5, 6},
    {1,4,11, 6}, {14, 1, 4, 9}, {4, 1, 14, 9}, {7, 1, 3, 10}, {3, 1, 7, 10}}
In[90]:= Length[S]
Out[90]= 6
```


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For a given family of graphs (paths, cycles, etc) are there better ways to count?

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## Lemma

If $v$ is a vertex of degree two which is connected to $u$ and $w$, and $d \mid r_{u}$ and $d \mid r_{v}$ then $d \mid r_{w}$.


Proof: $r_{v}\left|\left(r_{u}+r_{w}\right) \Rightarrow d\right|\left(r_{u}+r_{w}\right) \Rightarrow d \mid r_{w}$

## Lemma

If $v$ is a vertex of degree two which is connected to $u$ and $w$, and $r_{v}>r_{u}$ and $r_{v}>r_{w}$ then $r_{v}=r_{u}+r_{w}$.


Proof: $r_{v} \mid\left(r_{u}+r_{w}\right)<2 r_{v} \Rightarrow r_{v}=r_{u}+r_{w}$.

## Lemma

If $v$ is a vertex of degree two which is connected to $u$ and $w$, and $r_{v}>r_{u}$ and $r_{v}>r_{w}$ then $r_{v}=r_{u}+r_{w}$.


Proof: $r_{v} \mid\left(r_{u}+r_{w}\right)<2 r_{v} \Rightarrow r_{v}=r_{u}+r_{w}$.
Moreover, removing vertex $v$ gives a valid arithmetical structure on a (smaller) graph:


Let $P_{n}$ denote the path on $n$ vertices.

$$
r_{1}-r_{2}-r_{3}-\cdots-r_{n}
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Note: $r_{1} \mid r_{2}$, so by lemma $r_{1} \mid r_{3}$ and all other $r_{i}$.

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Note: $r_{1} \mid r_{2}$, so by lemma $r_{1} \mid r_{3}$ and all other $r_{i}$. But $\operatorname{gcd}\left(r_{1}, \ldots, r_{n}\right)=1$, so we must have $r_{1}=1$.

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Note: $r_{1} \mid r_{2}$, so by lemma $r_{1} \mid r_{3}$ and all other $r_{i}$.
But $\operatorname{gcd}\left(r_{1}, \ldots, r_{n}\right)=1$, so we must have $r_{1}=1$. Similarly, $r_{n}=1$.
$1-r_{2}-r_{3}-\cdots-1$


Note: No two consecutive entries are equal unless they are 1 .


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We can repeat this process until we get all 1's.

## Conversely...

Any arithmetical structure on $P_{n}$ can be obtained by starting with the trivial structure on $P_{k}$ for some $k \leq n$ and repeatedly adding vertices that are the sum of two adjacent vertices.

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$$
\begin{aligned}
& \text { Example }(n=4) \\
& 1-1-1-1
\end{aligned}
$$

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Analagous to some work of Richard Stanley, this implies that the number of structures is $1,2,5,14,42, \ldots$

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## Theorem

The number of arithmetical structures on $P_{n}$ is given by the Catalan number

$$
C_{n-1}=\frac{1}{n}\binom{2 n-2}{n-1}
$$

## Arithmetical Structures on Cycles

Let $C_{n}$ be a cycle on $n$ vertices.


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Let $C_{n}$ be a cycle on $n$ vertices.


Similar story works. Any two consecutive elements must be relatively prime. So either all entries are equal to one or there is a local maximum that is the sum of its neighbors.

## Arithmetical Structures on Cycles

Any structure can be attained by putting the number one in a positive number of vertices, and then filling others in with the sum of their neighbors.


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## Theorem

The number of arithmetical structures on $C_{n}$ with $k$ ones is

$$
\left(\binom{n}{n-k}\right)=\binom{2 n-k-1}{n-k}
$$

In particular, the total number of arithmetical structures on $C_{n}$ is

$$
\sum_{k=1}^{n}\left(\binom{n}{n-k}\right)=\left(\binom{n+1}{n-1}\right)=\binom{2 n-1}{n-1}
$$

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Let $Q_{n}$ be the path on $n$ vertices with first edge doubled:


The number of Arithmetical Structures on $Q_{n}$ is $4 C_{n-1}-2 C_{n-2}$.

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Let $Q_{n}$ be the path on $n$ vertices with first edge doubled:


The number of Arithmetical Structures on $Q_{n}$ is $4 C_{n-1}-2 C_{n-2}$. If first edge is tripled, there are:

$$
9 C_{n-1}-9 C_{n-2}-\frac{9}{n-1}\binom{2 n-6}{n-2}-\frac{4}{n-1}\binom{2 n-7}{n-2}-\frac{5}{n-1}\binom{2 n-8}{n-2}-\frac{6}{n-1}\binom{2 n-9}{n-2}-\frac{7}{n-1}\binom{2 n-10}{n-2}
$$

## Approach With Josh Wagner (Gettysburg '19)

Instead of counting the total number of structures, count the number of 'smooth structures'.

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Equivalently, no label $r_{v}$ is the sum of its neighbors.

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Instead of counting the total number of structures, count the number of 'smooth structures'.

## Definition

A smooth arithmetical structure is one where all $d_{i}>1$.
Equivalently, no label $r_{v}$ is the sum of its neighbors.
Seems as though it is easier to count these and then separately count the ways they can be subdivided to give all structures.

## Example

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$2 \backsim 3$

$2-1$

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- The final one can be subdivided in $C_{n-1}-2 C_{n-2}$ ways.
- Hence, total number of structures on $Q_{n}$ is $4 C_{n-1}-2 C_{n-2}$.

This approach works for other families of graphs

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$\square$

- Each of the first three can be subdivided to a structure on $Q_{n}$ in $C_{n-1}$ ways.
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- Hence, total number of structures on $Q_{n}$ is $4 C_{n-1}-2 C_{n-2}$.

This approach works for other families of graphs but appears to get very ugly very quickly.

## Reasons I like this topic

- Intermingling of additive number theory and multiplicative number theory.
- Gives a way to introduce topics in modular arithmetic while working on open research problems.
- Connects number theory to linear algebra and algebraic geometry.
- Very accessible to both computer experimentation and exploration 'by hand'.
- Lots of fruit that appears to be low-hanging...


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